

Hop-by-hop Adaptive Video Streaming in Content Centric Network

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Background

- Background:
 - **Content-Centric Networking (CCN)**
Future network proposals
Focuses on content retrieval with content-based addressing
Universal caching design
 - **Adaptive video streaming (DASH)**
Dynamically adjust the bitrates for video services
 - **Scalable video coding (SVC)**
Encode video into base and enhancement layer content
Support progressive transmission

Aim

- Problems for DASH over CCN:
 - Bitrate oscillations
 - Network congestions

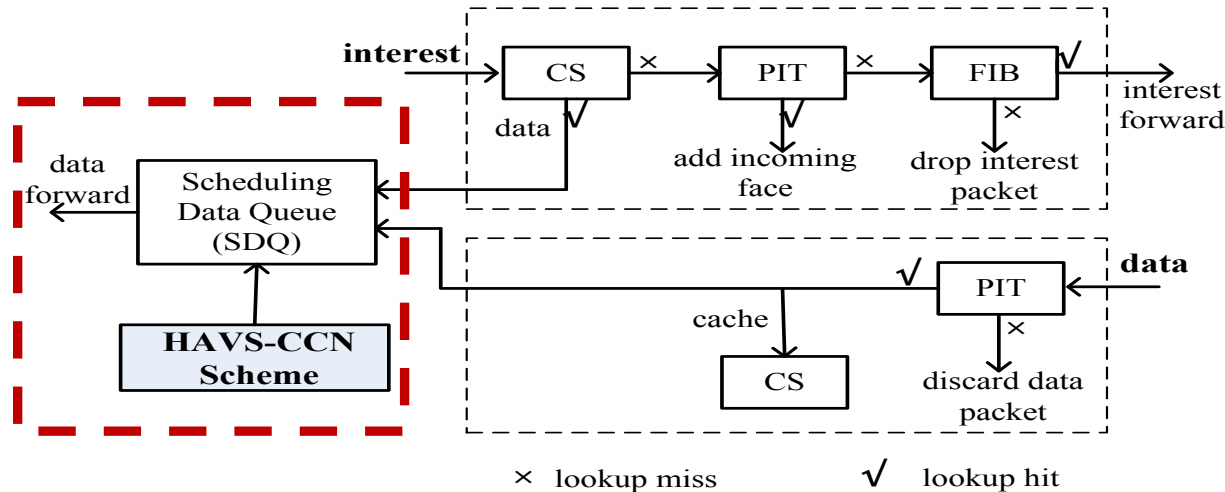
We propose a hop-by-hop adaptive video streaming scheme (**HAVS-CCN**)

- Solved by:
 - Adjust the video quality hop-by-hop
 - A flow control for data packets, directly

SYSTEM MODEL

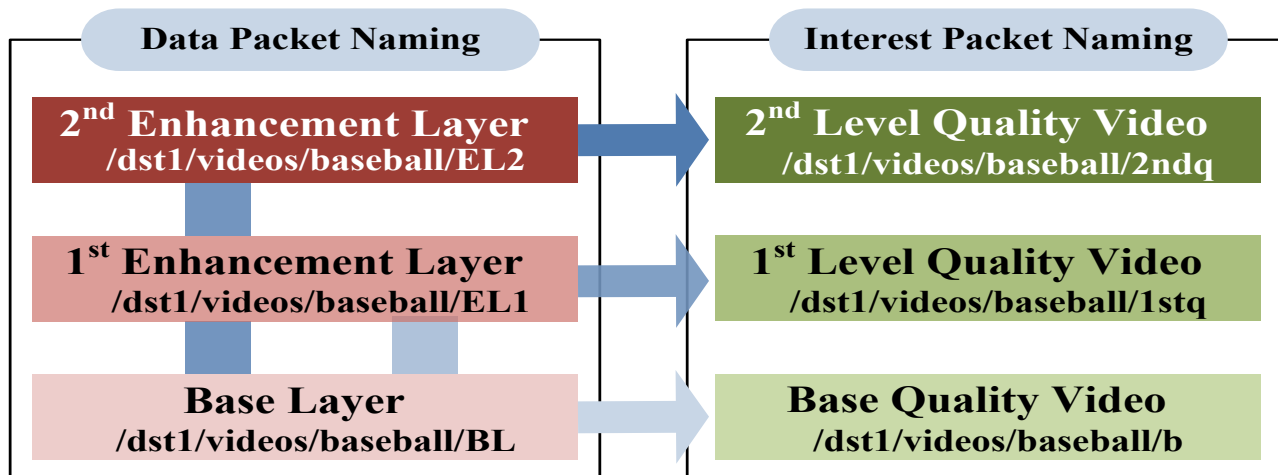
System model

- Improved CCN router design
 - Scheduling Data Queue (SDQ)
 - Schedule the outbound data flow for each interface



System model

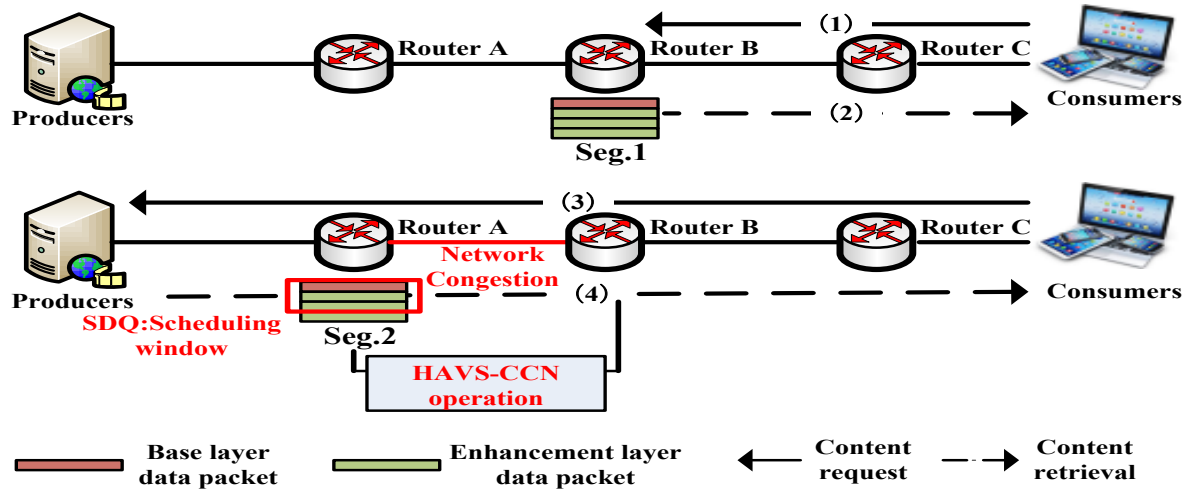
- Scalable Video Segmentation and Naming



CCN nodes can directly drop the some enhancement layer packets to guarantee **smooth playback**.

Building HAVS-CCN

HAVS-CCN overview



1. Set a scheduling window

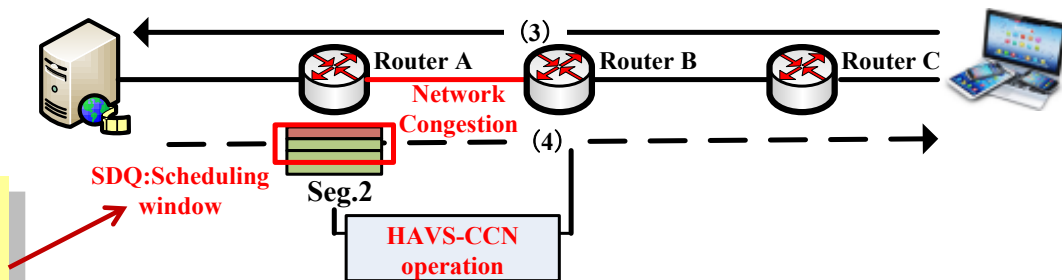
2. Compute priority for SVC data packet

3. Hop-by-hop operation

Description of scheme (1)

- Set a scheduling window to limit the video transmission rate
 - Scheduling Data Queue (SDQ)
 - Schedule the outbound data flow for each interface

$$wnd_{scheduling}(k) = delay \cdot \alpha \cdot BW_{out}(k)$$



Scheduling window to make flow control

Description of scheme (2)

- Priority of each SVC content data packet

Type-based priority

$$w_t(p_{i,seg,type}) = \frac{1}{\eta_{type}}$$

$$\eta_{type} = \begin{cases} 1, & type = base\ layer \\ i + 1, & type = i^{th}\ enhancement\ layer \end{cases}$$

Popularity-based priority

$$w_p(p_{i,seg,type}, k) = \frac{\eta_{i,k}}{SDQ_k}$$

- A data packet with higher **type-based priority** and higher **popularity-based priority** will have preference to be transmitted

Overall priority: $w(p_{i,seg,type}, k) = w_t(p_{i,seg,type}) \cdot w_p(p_{i,seg,type}, k)$

Description of scheme (3)

- The aim of the HAVS-CCN is to transmit high priority data packets first when the link capacity is limited

Objective:

$$\text{maximize } f = \sum_{i, \text{seg}, \text{type}} c_{i, \text{seg}, \text{type}} \cdot e^{w(p_{i, \text{seg}, \text{type}}, k)}$$

Subject to:

$$\begin{aligned} \sum c_{i, \text{seg}, \text{type}} \cdot p_{\text{size}}(\text{type}) &= \text{wnd}_{\text{scheduling}}(k) \\ c_{i, \text{seg}, \text{type}} &\in \{0, 1\} \\ c_{i, \text{seg}, \text{type}} &\neq 0, \text{type} = \text{base} \end{aligned}$$

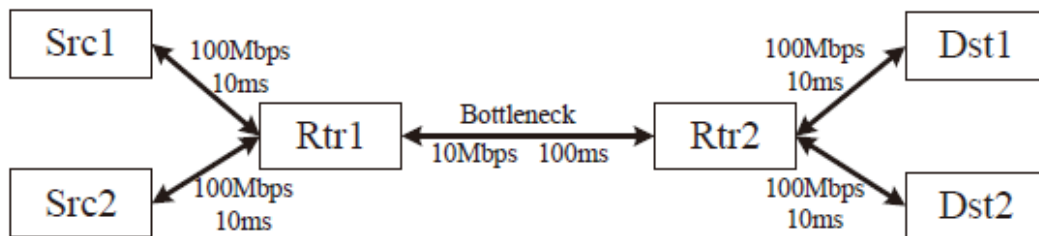
Algorithm 1 HAVS-CCN Priority Transmission Algorithm

```
1: procedure ONWINDOWEMPTY(interface)
2:   TxQueue  $\leftarrow$  SDQ[interface].
3:   if sizeof(TxQueue) > wndscheduling then
4:      $c_{i, \text{seg}, \text{type}} \leftarrow$  Solver( $p_{i, \text{seg}, \text{type}}, w_{i, \text{seg}, \text{type}}, k$ )
5:     TxQueue  $\leftarrow$  { $p_{i, \text{seg}, \text{type}} \mid c_{i, \text{seg}, \text{type}} \neq 0$ }
6:   end if
7:   Send(TxQueue).
8: end procedure
```

Simulation

Small scale simulation

- 6-node topology
- Demonstrate the real-time response of HAVS-CCN when the available bandwidth changes.

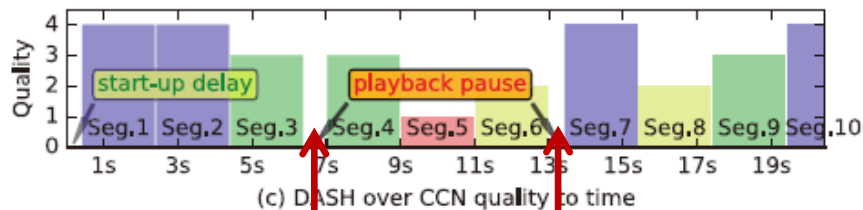
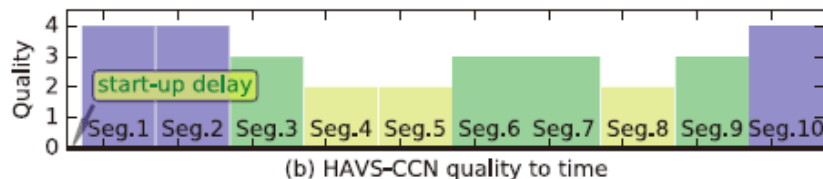
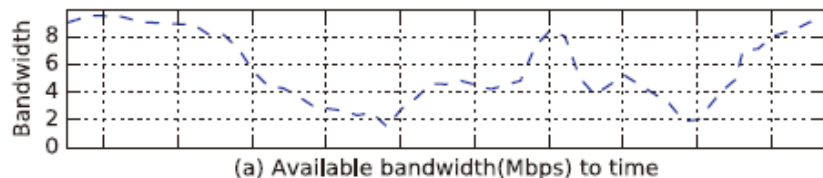


- Video source: SVC encoded video from University Klagenfurt [1]
- The demo video has 500 frames which are encoded into 11 segments. Each segments have a base layer and 3 enhancement layers.

[1] C. Kreuzberger et.al, "A scalable video coding dataset and toolchain for dynamic adaptive streaming over http," in MMSys 2015

Small scale simulation

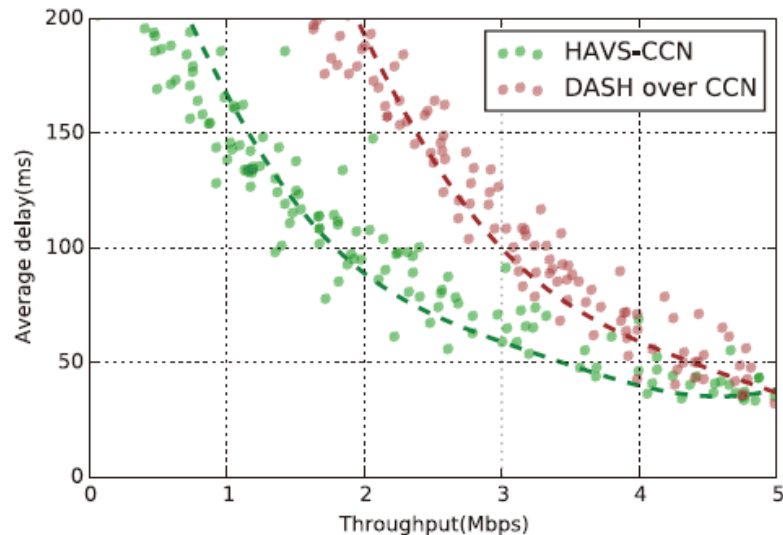
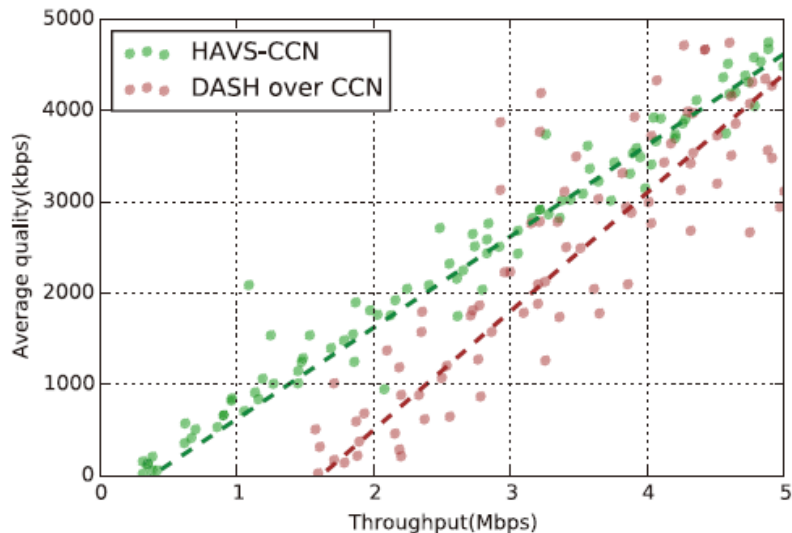
- The aim of the HAVS-CCN is to transmit high priority data packets first when the link capacity is limited



- HAVS-CCN reduce playback pause:
1. when the link encounters a varying **pressure** from *Src2*
 2. when **bandwidth overestimation**, can directly drop some *enhancement layer packets*

Large scale simulation

- Realistic topology: Rockfuel's AT&T topology [1], including 192 nodes classified into clients, gateways and backbones



Metrics: average quality and average delay

[1] N. Spring, R. Mahajan, and D. Wetherall, "Measuring isp topologies with rocketfuel," in *ACM SIGCOMM Computer Communication Review*

Conclusion

- We introduce a hop-by-hop solution which achieves:
 - reduce the negative effect of inaccurate bitrate predication
 - provide flow controls at router level
- Novelty:
 - An novel add-on, SDQ component, queues data packets for each interface.
 - The video quality can be adjusted directly along the data packet transmission path

Future work

- Mathematically analyze the priority weights setting for SVC packets

Questions, discussion...

- Thank you very much for your kind attention!
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